

## I. IP Fragmentation (pg 45-52)

Week 9

- recall IP is responsible for 1) network addressing and 2) fragmentation
- the 2<sup>nd</sup> “word” of the IP header is the fragmentation related data, composed of the **ID** (16 bits), the **fragment flags** (3 bits), and the **fragment offset** (13 bits) (see IP PDU)
- physical layer typically imposes an upper limit on the size of a frame, e.g. MTU (max Ethernet frame data size?)
- IP compares the datagram size (as stored in the ? field) with the MTU
  - if  $TL > MTU$  then  
fragmentation should occur
- fragmentation may occur on original host, or on intermediate nodes (e.g. routers) along the way
- datagrams & fragments can be fragmented more than once
- once fragmentation occurs, reassembly does not take place until the final destination
- each datagram fragment becomes an individual packet, so gets its own:
  - header (+ 20 octets per fragment)
  - data
  - route to destination and arrival order
- only data gets fragmented, never the header
- if one fragment is lost, the entire datagram is useless and must be retransmitted
- IP datagram – IP header and all data (i.e. all fragments)
- IP packet – IP header along with some data
- on Windows from cmd prompt: **netsh interface ipv4 show interfaces**

## How does fragmentation work: (see IP RFC link on syllabus)

1. IP layer queries the MTU of physical layer
  - if  $TL \leq MTU$  then
    - transmit datagram
  - elseif DF flag = 1 then
    - drop datagram
  - else
    - fragment datagram
2. number of fragments  $NF = (TL - IHL) / (MTU - IHL)$  (taking the ceiling of this)
  - note: IHL is 5 “words” \* 4 octets/word = 20 octets
3. # fragment blocks (8 octets per block)  $NFB = INT((MTU - IHL \text{ octets}) / 8)$
4. copy original header to the first (eventually all) fragment headers, including the same ID, protocol, destination and source addresses
5. fill in data and update header fields:
  - append  $NFB * 8$  data octets, unless last fragment
  - $TL = NFB * 8 + IHL$  octets, unless last fragment
  - Fragment Offset (FO) = (fragment # - 1) \* NFB
  - set may frag flag = 0
  - set more frag flag = 1, unless last fragment, then set = 0
  - decrement TTL if at router
  - recalculate checksum
6. transmit fragment
7. goto # 4 for next fragment

**Fragmentation Example** (see pg 49)

- beginning total length: 4464 octets, fragment?
- MTU: 1500

original datagram:

<b>Frag #</b>	<b>Frag ID</b>	<b>DF</b>	<b>MF</b>	<b>Frag Offset</b>	<b>IHL</b>	<b>TL</b>	<b>data len</b>
-	321	0	0	0	20	4464	4444

fragmented datagram (packets):

$NF = (TL - IHL) / (MTU - IHL) = 4444/1480 = 3.0027$ , take next largest INT value, so 4 fragments

$NFB = ((MTU - IHL) / 8 \text{ octets/block}) = (1500 - 20) / 8 = 185$

<b>Frag #</b>	<b>Frag ID</b>	<b>DF</b>	<b>MF</b>	<b>Frag Offset</b>	<b>IHL</b>	<b>TL</b>	<b>data len</b>
1	321	0	1	0	20	1500	1480
2	321	0	1	185	20	1500	1480
3	321	0	1	370	20	1500	1480
4	321	0	0	555	20	24	4

└───┬───> 4524 4444

Note that the sum of the data lengths = sum of original data length, but the sum of the total lengths is original TL + ((NF -1) \* 20) for extra headers