

I. Media Contention & MAC Rules

Week 5

Collisions in Ethernet (specifically 10Mbps – half duplex):

- are not errors
- do not corrupt data
- are the basis for multiple node access on a single half duplex line

10baseT manages collisions using something called **CSMA/CD** (Carrier Sense Multiple Access/Collision Detect).

CSMA/CD (see Ch 2 & Ch 3, pg. 48+)

1. node wishing to transmit listens to media for presence of a carrier signal (i.e. a signal on the wire)
 - if a carrier is present, node must wait
 - if no carrier for IFG time period, transmit
2. if a collision is detected, a (32 bit) **jam signal** is transmitted
3. after sending a jam, all nodes will wait a random amount of time before trying to send again (backoff)
4. goto step 1

Timing Issues:

- **network diameter** – distance between two (most distant) end nodes
- **network segment**
 - 1) in general, network section that is bounded by bridges, routers, or switches (i.e. a collision domain) per Cisco definition
 - 2) physical wire connection w.r.t. 5-4-3 rule, not a collision domain
- while signals on the line are very fast (almost 3×10^8 m/sec), they still take measurable time. Time is measured in **bit time**, for 10 Mbps:

10 Mbps = 10 million bits/sec or 10 bits/ μ second, so 1 bit 1/10 millionth of a second or 1 bit transmitted in .1 μ second
- **interframe gap** (IFG) – time for brief recovery after receiving (tail end of) frame
 - 96 bit times
 - 9.6 μ seconds (microseconds) for 10 Mbps

- **slot time** – maximum time required to detect a collision
 - smallest Ethernet frame? 64 octets, or 512 bits, thus 51.2 μseconds
 - approximately twice the signal propagation time between the two end nodes (most distant) on the network
 - why twice? the time it takes a signal to travel from one end to the other end, and back → **physical layer round trip propagation time**
 - slot time is the maximum time required for collision enforcement

- **wire speed & data delivery** (see Gallo, pg 334)
 - smallest Ethernet frame (worst case scenario) transmitted in 51.2 μsec + preamble 6.4 μsec +IFG 9.6 μsec totals **67.2 μsec per frame**
 - thus 1 frame is transmitted in $1/67.2 \mu\text{sec} = 0.0148809 \mu\text{sec}$
 - since 1 second is $10^6 \mu\text{sec}$, $0.0148809 \mu\text{sec} * 10^6 = 14880$ (min size) frames are transmitted in 1 second → this is referred to as **wire speed**
 - **data (payload) delivery:** at 14880 frames per second, 46 data octets (min) per frame, 8 bits/octet, we get $14880 \text{ frames/sec} * 46 \text{ octets/frame} * 8 \text{ bits/octet} = 5.476$ (approx) Mbps data

 - largest Ethernet frame: $1518 \text{ octets} * 8 = 12144 \text{ bits}$ transmitted in $1214.4 \mu\text{sec}$ per frame, + 6.4 + 9.6 = 1230.4 μsec
 - 1 frame in $1/1230.4 \mu\text{sec} = 0.0008127 \mu\text{sec}$ (approx)
 - $0.0008127 * 10^6 = 812.74$ frames per second
 - 1500 data octets (max) per frame, 8 bits/octet → $812.74 * 1500 * 8 = 9.773$ Mbps data

II. CSMA/CD & Backoff Algorithm

- goal: once a collision occurs, goal is to minimize chances of another
- backoff algorithm provides a mechanism for adjusting timing of retransmissions
- based upon slot time multiple with random integers in a given range
- integer (i) chosen such that $0 \leq i \leq 2^n - 1$ where n is collision attempt and $n \leq 10$

collision #	range	wait time
0	-	-
1	0-1	0 * 51.2 μ s, 1 * 51.2 μ s
2	0-3	0 * 51.2 μ s, 1 * 51.2 μ s, 2 * 51.2 μ s, 3 * 51.2 μ s
10	0-1023	
11-15	same as above	
16	discard frame	

III. Ethernet Restrictions

10 Mbps Ethernet:

	Type	Max Seg Len	Max Nodes Seg	Max Diameter	Spacing
Thick coax	10base5	500 m	100	2500 m	2.5 m
Thin coax	10base2	185 m	30	925 m	.5 m
TP	10baseT	100 m	2	500 m	-
fiber	10baseF	2000 m	2	2500 m	-

- see hub handouts and network diameters

IV. 5-4-3 Rule

- rule is applied to Ethernet networks where greater distances are achieved using (repeating) hubs
- used to insure that any data sent will arrive at its destination within a given time, thus reducing collisions
- only applied to hub based Ethernets, does not apply to switched Ethernets
- rule states:
 - given two types of segments, user (populated) segments and link segments
 - for any two end nodes on the LAN, there can be a maximum of **5 segments** between them
 - connected via a maximum of **4 (repeating) hubs**
 - where a maximum of 3 **of the segments are user populated** (i.e. have user nodes connected to them)
 - sometimes called 5-4-3-2-1, where 2 is the number of segments that are not user populated and 1 is the number of collision domains
- workarounds to 5-4-3, stackable hubs, but these are also limited based upon manufacturer
- see handout